# Peer to Peer VPNs

### Enterprise Networks (6CFU)

Slides taken from Prof. Andrea Detti's TPI2 lectures.

## Introduction

- Allow secure traffic exchange among company branches ditributed over the entire territory
- Usually required by business customers
- Virtual Private Networks
  - » Private: allows communication between subnets in different networks as they were in the same private network (as for addressing, routing and security)
  - » Virtual: the required links between networks are (necessarily) virutal (not phisical). The support network is not private.

### • Private IP addressing

- » 10.0.0/8
- » 172.16.0.0/12
- » **192.168.0.0/16**

### • **Requirements**: unique addressing in a VPN

# **VPN models**

### Communication models

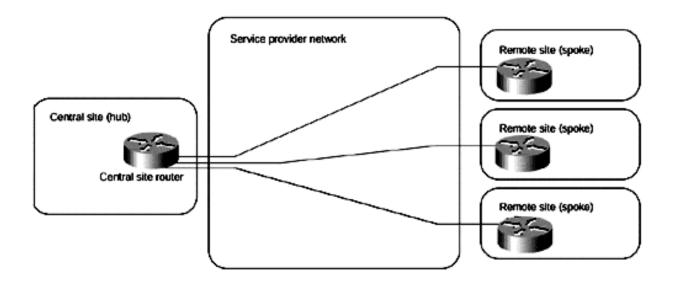
- » Intra-company (Intranet)
- » Inter-company (Extranet)
  - » Addresses must be unique
- » VPDN (Virtual Private Dialup Network)
  - » Dynamic address configuration

### Data transfer models

- » Overlay: ISP network is used only for transporting features. Routing information is exchanged between company networks. The VPN topology composed by point to point links configured by the customers
- » Peer-to-peer: the ISP is responsible also for exchangin routing information Logical topology is defined by the customers. Physical topology is defined by the ISP

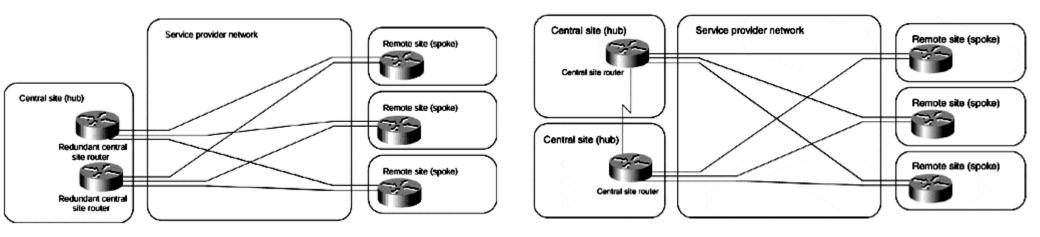
### **VPN topologies – Hub and Spoke**

- VPN topology depends on the specific customers' needs. Nevertheless, there are some "standard" topologies...
- Hub-and-spoke Topology:
  - » Remote branches (spoke) connected to a central site (hub).
  - » Spokes can communicate with each other, but inter-spoke should be negligible then spoke-hub traffic



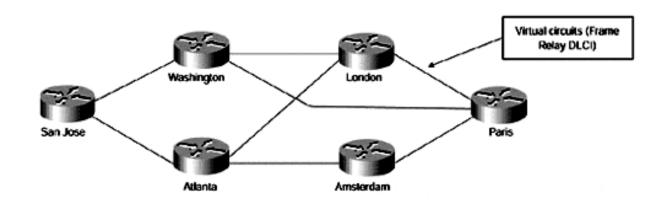
### **VPN topologies – Hub and Spoke**

• Hub Backup



## **VPN topologies – Partial/Full-Mesh**

- When there's a huge data exchange between enterprises' sites, the Hub-and-Spoke topology is less effective since all the spoketo-spoke traffic traverses the hub → bottleneck
- In such a case, partially or totally connected topologies are preferred
- Business case:
  - » Companies without a strict gerarchic organization
  - » Peer-to-peer applications (messaging or collaboration system)
  - » For multinational companies in which the cost of the hub-and-spoke solution could be high because of the excessive cost of international links.

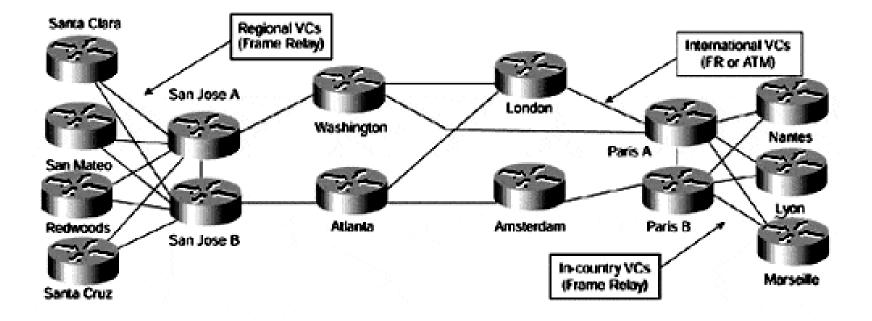


## **VPN topologies – Partial/Full-Mesh**

- The full-mesh topology is simple to plan
  - » Get the traffic matrix A(i,j)=x Mb and ask to the ISP a link between site i and site j with x Mbps
- BUT... the full-mesh cost can be high since the number of links employed is n\*(n-1)
- SO... often a partial mesh topology is adopted
- How to plan a partial mesh topology?
  - » 1) Create a fully connected topology through links only between sites having a huge traffic exchange
  - » 2) From the traffic matrix, and assuming a shortest-path routing, compute the amount of required bandwidth on all the employed links
  - » 3) Order links from the cheapest ISP ;-)

### **VPN topologies – Hybrid**

- Huge international VPNs are often composed by many national hub-and-spoke VPNs
- The international part (backbone) is a partial-mesh between hubs



## **Peer-to-Peer VPN**

### **Peer-to-Peer VPN**

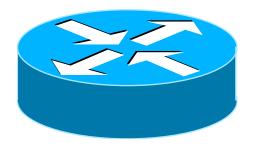
- Routing Information exchange between Company and ISP routers → routing happens on a layer composed both by company entities and by ISP entities
- De facto based on BGP/MPLS solution
  - » Enterprise's gateway transfers data to the ISP which handles the forwarding through other Enterprise's sites
  - » Routing (connections topology) is actually in the hands of the ISP
  - » Plug & Play, adding a site is a matter of ISP configuration only, the company has to do almost nothing

# **VPN BGP/MPLS**

**Peer-to-Peer VPN** 

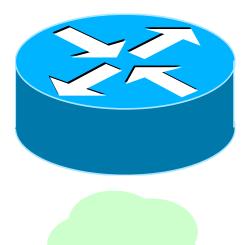
# Elements of a VPN BGP/MPLS network





**Customer Edge** : is the Company side router facing with the ISP which provides the VPN BGP/MPLS service. It has standard routing functionalities; its only peer is the Provider edge with which exchanges info through BGP messages

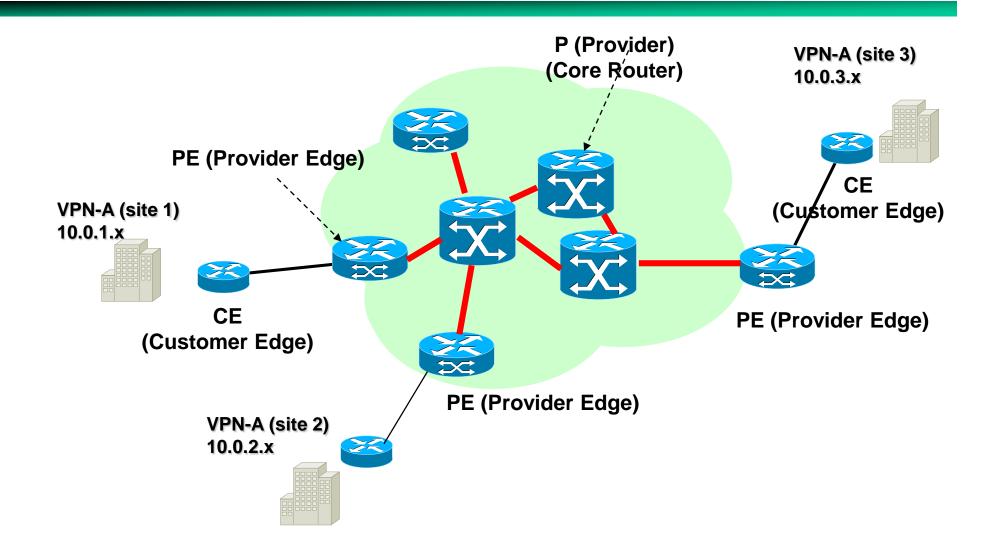
**Provider Edge** : is the access router on the ISP side in which one or more Customer Edges are connected. Besides IP functionalities, it also handles the MPLS LER role.



**Provider Router** : Label Switched Router (LSR) composing the MPLS backbone of the ISP

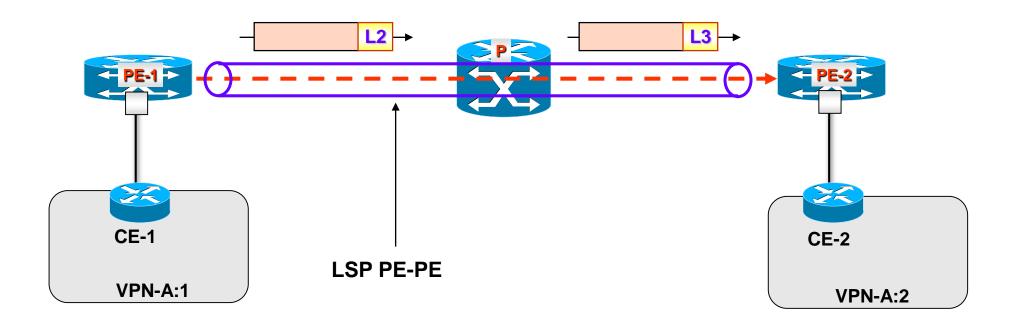
**MPLS/VPN Backbone** : MPLS network with properly configures LSPs to interconnect all the Provider Edges.

# **VPN MPLS service architecture**

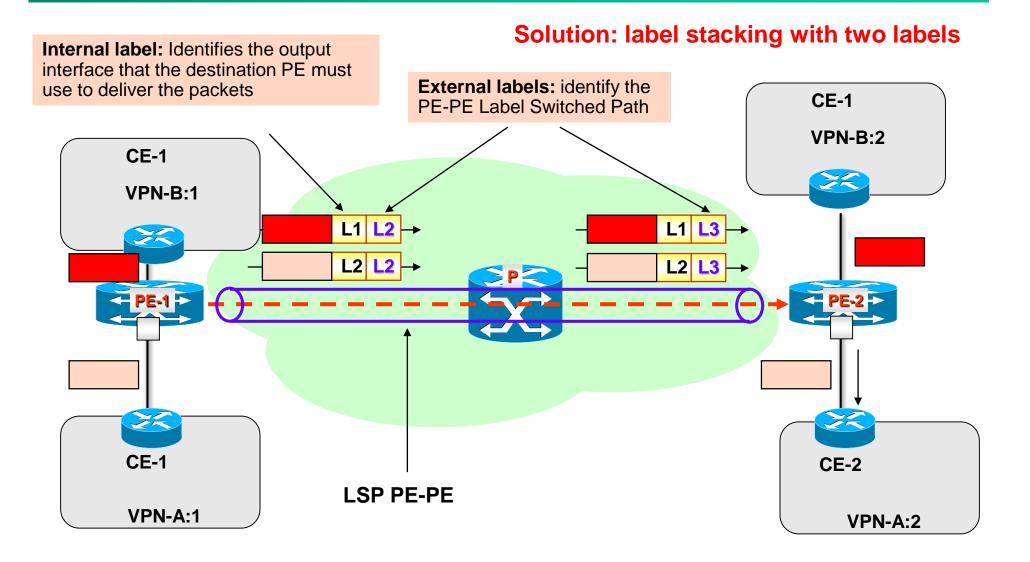


# **Forwarding mechanism**

- Problem: transfer packets between the two sites of a VPN → A:site1 --- A:site2
- Trivial solution (A:1 → A:2): encapsulate at PE (A:1) IP packets coming from CE(A:1) in the ISP, linking PE(A:1)→PE(A:2)
- At the LSP end, PE(A:2) forwards on an IP base
- What happens if same PEs support more than a VPN with *non-coordinated* addressing?
- It can happen that PE(A:2) finds itself to forward (on IP base) packets with the same network addesses, BUT different VPNs!!!



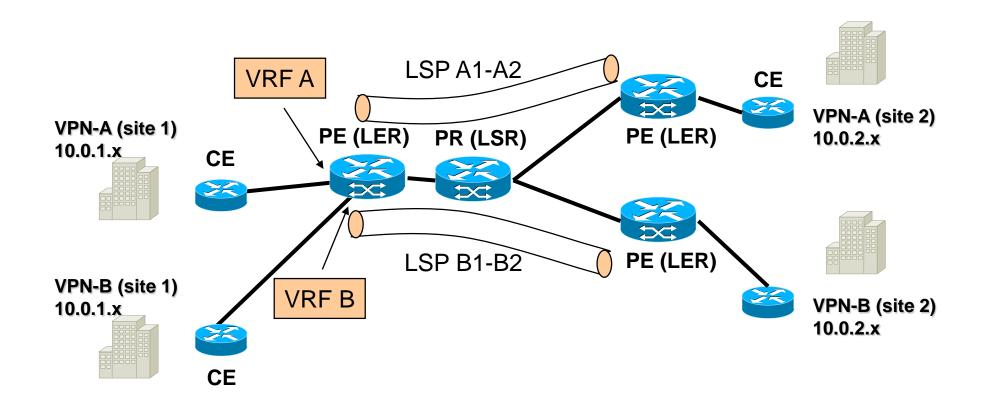
# **Forwarding mechanism**



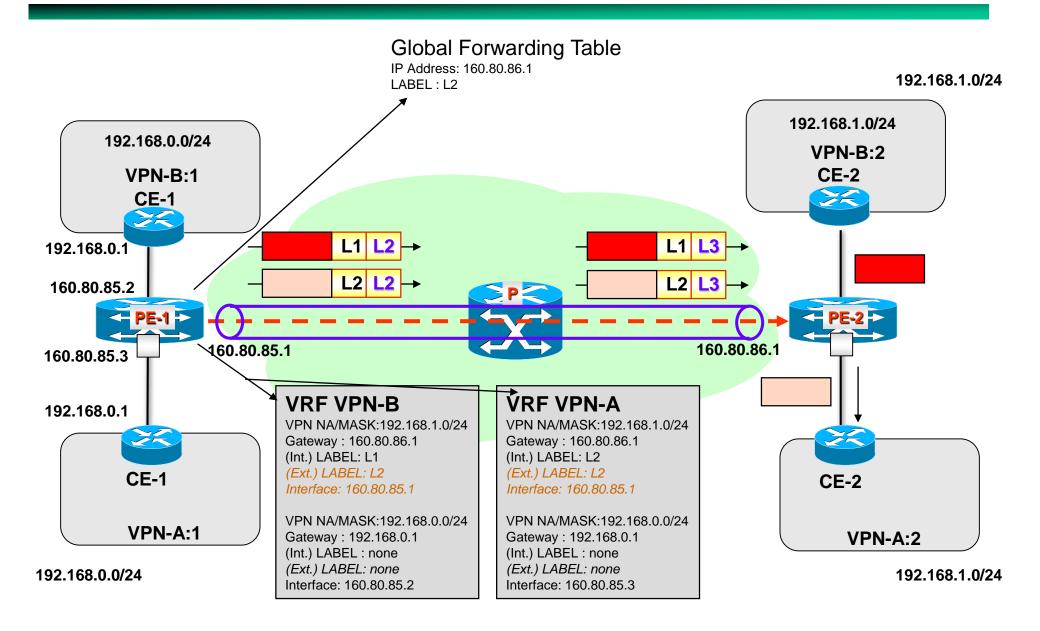
### **PE classification**

- PROBLEM: how can the PE forward/classify the packets coming from CE A:1 on the right tunnel?
- **SOLUTION**: it must know to which VPN packets belong to
  - » this information is deduced from the ingress interface in which a packet is received
- SO, depending on the belonging VPN, the MPLS forwarding of the packet changes. Technically, the PE stores as many forwarding tables as the number of VPNs connected to it. Each virtual table is named VPN Routing and Forwarding (VRF) table
  - » A VRF entry contains (logically) the following tuple: <VPN network address, VPN mask, Next PE IP Address, Internal label, Output Interface>
- In addition to the VRF, a PE stores one Global Forwarding Table (GRT) which permits to reach a PE from another PE.
  - » Logically, a GRT entry contains the tuple: <PE IP address, external label, Output Interface>

### **PE classification**



# **VRF and GFT**



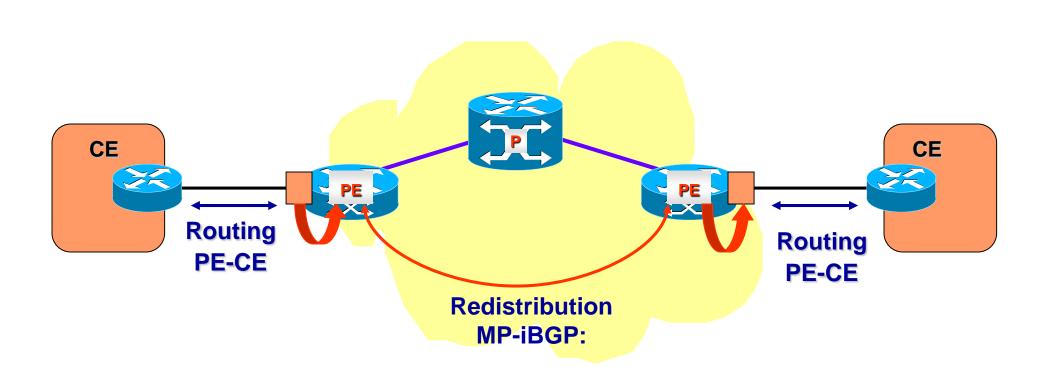
# **Populating GFT and VRFs**

- The Global Forwarding Table is configured by the provider during the set-up o the MPLS/VPN backbone (i.e. LSPs between PEs)
- The GFT can be populated manually (in the case of manual LSPs), or automatically in the case of a set-up with signalling protocols like LDP, RSVP-TE or CR-LDP
- VRFs contain two forwarding categories:
  - » Forwading to LOCAL sites
  - » Forwarding to REMOTE sites

### • Forwarding to local sites can be:

- » Manually configured
- » Obtained through specific routing protocols (OSPF, RIP, etc.), running the CE-PE link
- Remote routes are obtained through an extension of the BGP-4 protocol, namely Multi-Protocol interior BGP (MP-iBGP or MP-BGP)

# **Populating VRFs**



- Routing CE-PE: Static, RIP, OSPF, eBGP
- Routing PE-PE: MP-iBGP = MultiProtocol-internal BGP

## **Populating VRFs**

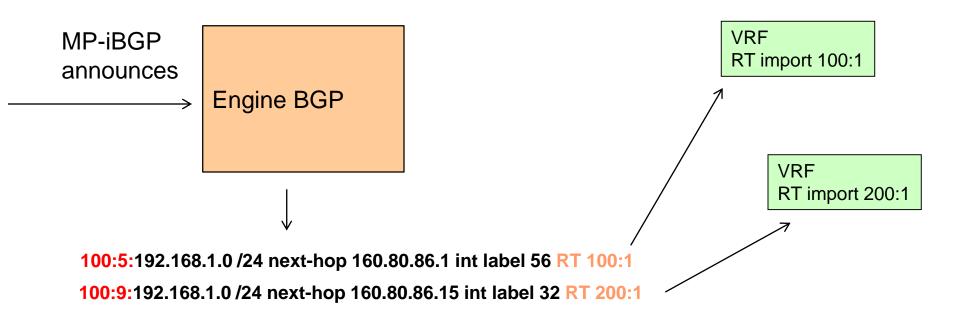
- VRFs "synchronize" among them exchanging the reachability info inside MP-iBGP announces
- An MP-iBGP announce is sent by a PE to all other PEs; i.e. it exists an overlay full mesh between PEs
- Assumption: the cost of the *direct hop* between two PEs is 1, being this an IP level hop (not MPLS hop)
- A same MP-iBGP announce carries reachability information relative to prefixes of more VRFs

## **Route Distinguisher**

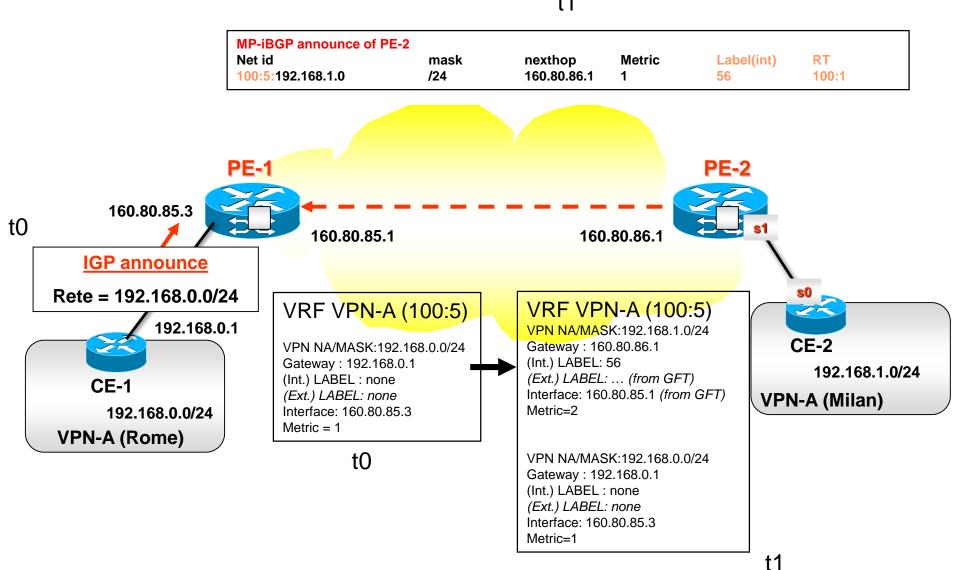
- Thanks to MP-iBGP announces, the BGP engine inside the PE calculates the next-hop (and internal label) towards every announced prefix
- VRFs belonging to different VPNs can notify a same private prefix since the addressing spaces can be overlapped.
- To differentiate overlapped prefixes (i.e. make them different to the BGP engine), a VRF is characterized by an ID named Route Distinguisher (64 bit)
  - » Usually, all the VRFs of the same VPN use the same Route Distinguisher, since the prefixes inside a VPN cannot overlap.
  - » In this way, the Route Distinguisher can be reused

### **Route Distinguisher**

- The RD is placed before the net\_id in the MP-iBGP entries
- The routes computed by BGP are inserted inside the enabled VRFs (see Route Target next...)



# **Populating VRFs**



#### t1

## What about the VPN topology?

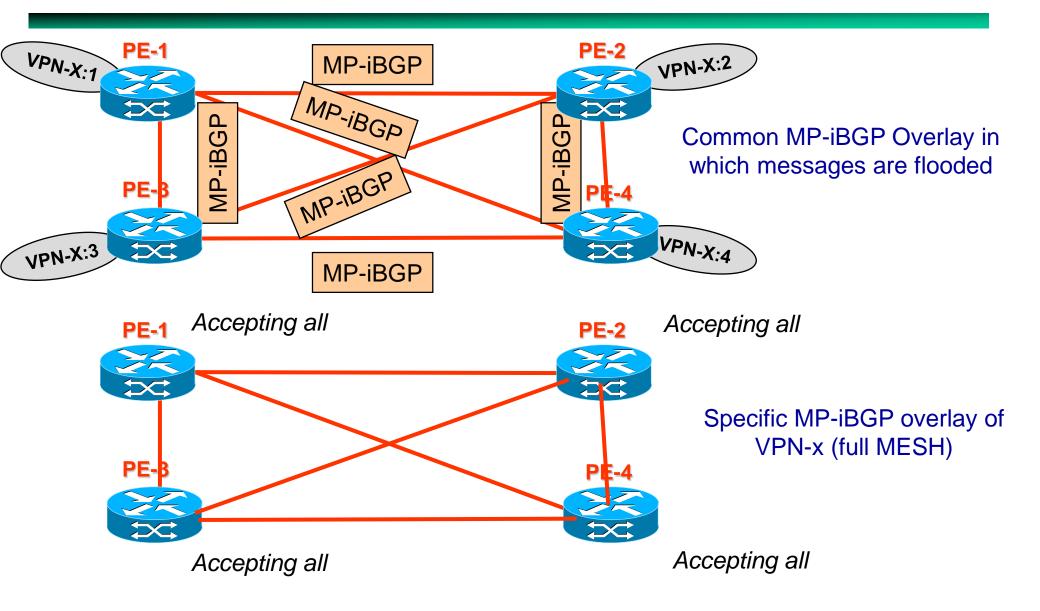
- If MP-iBGP messages are diffused among all PEs, all the VPNs have a full-mesh topology
- PROBLEM: what if I want different topologies for different VPNs?
- BGP priciples say that if I have an overlay topology in which MP-iBGP messages are diffused, the (forwarding) topology of VPN-x is the set of the overlay shortest-paths between any couple of nodes
- Since direct connections between two PEs have metric 1
   → the VPN-x topology matches the overlay topology in
   which MP-iBGP messages are notified
- Therefore, if the overlay network in which MP-iBGP messages are forwarded is full-mesh, the VPN topology is full-mesh, too

## What about the VPN topology?

 To change the logical topology of VPN-x it is necessary to change the MP-iBGP overlay network of VPN-x

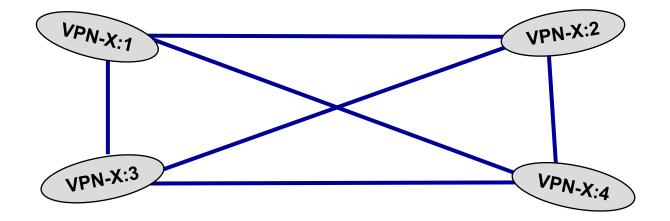
- » Solution 1: create a different MP-iBGP overlay forwarding topology for each VPN
  - » Cons: high management effort, cannot aggregate inside the same MP-iBGP message the routing information relative to more VPNs, etc.
- » Solution 2:
  - » Having an overlay full-mesh for MP-iBGP common between PEs
  - » Define the *specific* overlay needed for a given VPN
  - » Flood MP-iBGP messages on the common MP-iBGP overlay
  - » Receivers elaborate only announces coming from links of the specific overlay

### **Populating VRFs - VPN Full Mesh**

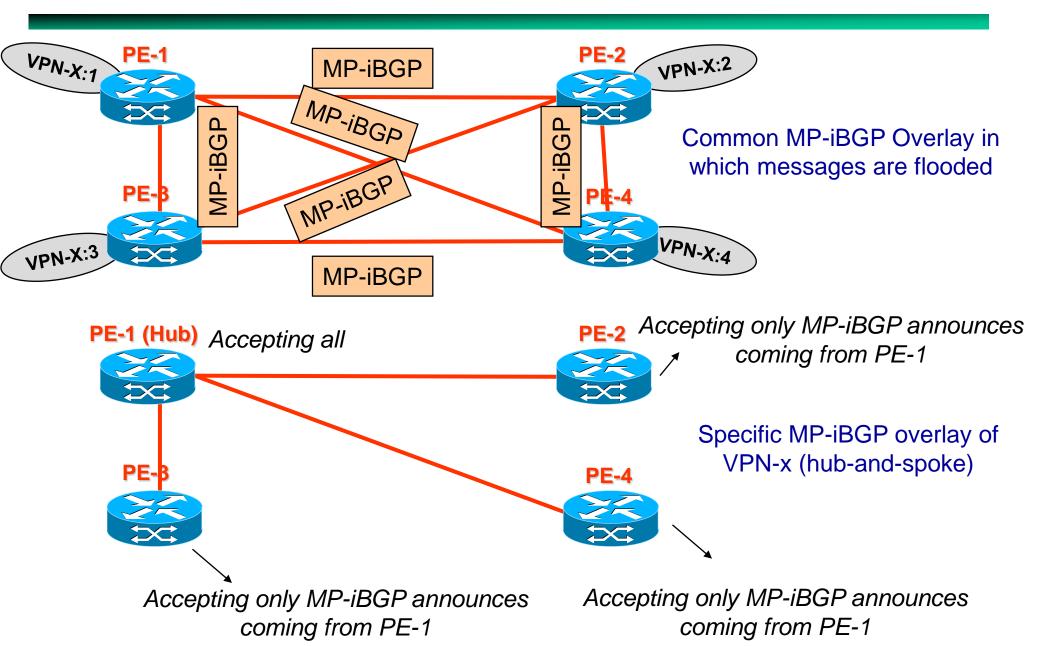


### **Populating VRFs - VPN Full Mesh**

Resulting VPN topology

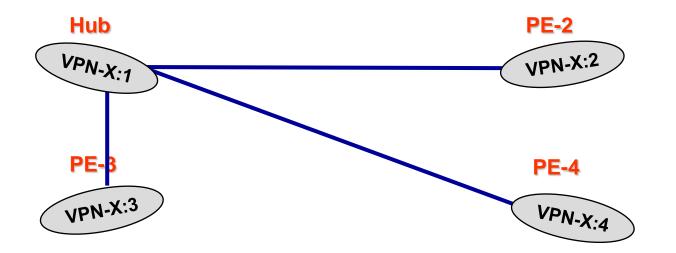


### **Populating VRFs - VPN Hub and Spoke**



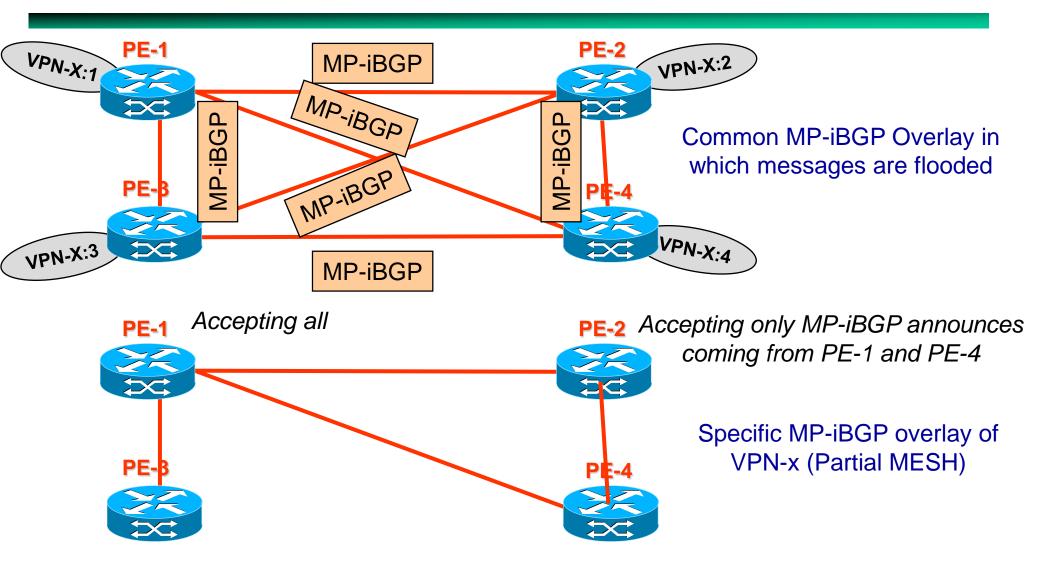
### **Populating VRFs - VPN Hub and Spoke**

### **Resulting VPN topology**



**NOTE**: MP-iBGP announces (like iBGP) cannot propagate on iBGP links. Therefore, in order to permit to spokes the communication between them, The hub's VRF must export a default route

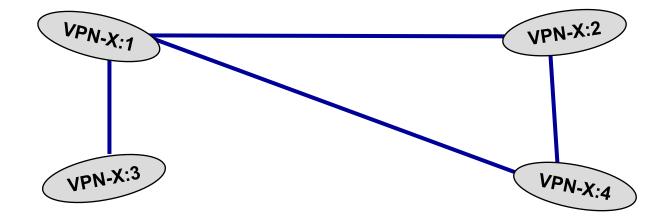
### **Populating VRFs - VPN Partial Mesh**



Accepting only MP-iBGP announces coming from PE-1 Accepting only MP-iBGP announces coming from PE-1 and PE-2

### **Populating VRFs** - VPN Partial Mesh

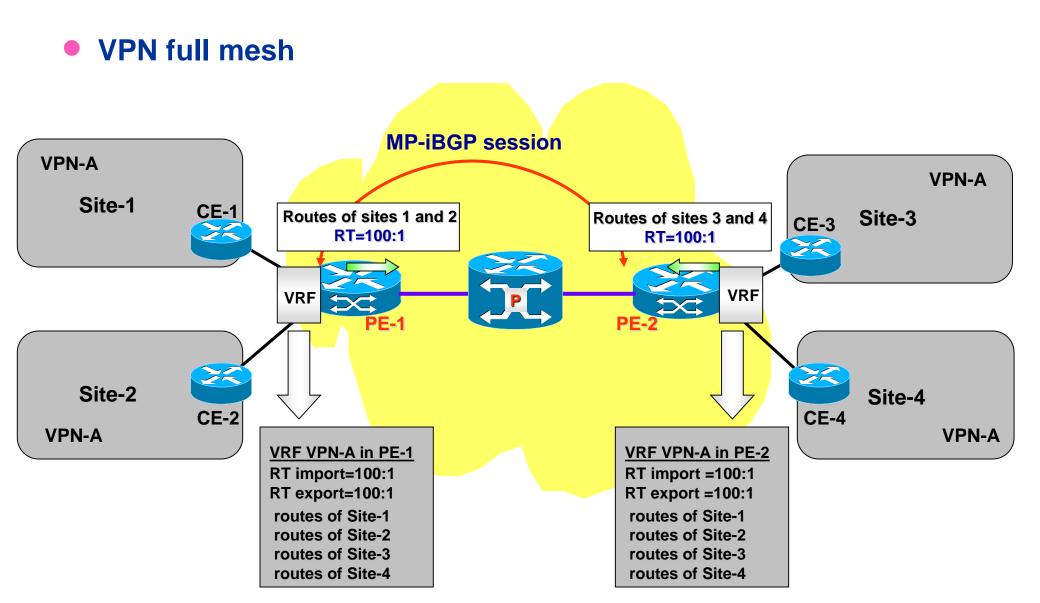
Resulting VPN topology



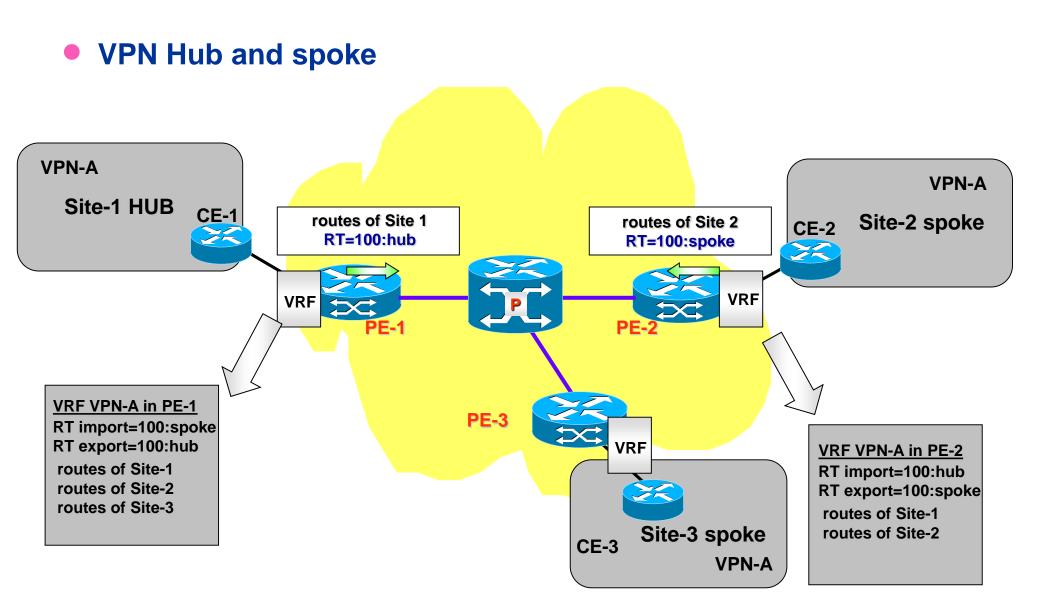
# **Route Target**

- <u>The Route Target concept</u> permits to realize a specific overlay for the VPN-x discussed before. Therefore, permits to define VPN-x topology.
- It's the VPN/MPLS "way" to tell to a VRF-x to "accept only a subset of MP-iBGP announces"
- HOW:
  - » Each VRF transmitting announces, labels (exports) these announces with a configurable ID (Route target) of 64 bit size
  - » Each VRF can receive (import) only announces with a configurable subset of Route Targets

### Using the "Route Target": Example 1



### Using the "Route Target": Example 2



# **VPN/MPLS** configuration

### • Initialization

- » Configure LSP MPLS (e.g. with LDP) between all PEs
- » Enable BGP peering for prefixes of type *VPNv4* (RD+net\_id) between all PEs
- How to add another site:

### Client

- Notify to ISP the need of another VPN site and the relative topology
- » Install a CE as enterprise gateway
- » Configure the *default gateway* of the CE with the IP address of the access PE
- > Optional: enable on CE a routing protocol on the CE-PE path (e.g. OSPF)

### Provider

- » Initialize a new VRF on access PE
- » Define/Configure the Route Distinguisher
- » Define/Configure Route Import and Route Export and eventually update the import/export RTs on the other PEs, coherently with the requested topology
- » Associate the ingress PE interface with the VRF
- » Enable MP-iBGP on such VRF

### **VPN/MPLS** conclusions

- Simple approach for the client
- Security of traffic is delegated to the ISP (not to be handled by the client!)
- Modest configuration complexity by the provider
- Provider must «traffic engineer» the VPN/MPLS backbone so to offer the promised QoS to the VPN clients
  - » Traffic engineering in the LSPs between Pes
- The cost could be high, though

### On all involved PE

- Create user VRF
  - » PE-1(config)# ip vrf vpnB
  - » PE-1(config-vrf)#rd 200:0
  - » PE-1(config-vrf)#route-target import 200:2
  - » PE-1(config-vrf)#route-target export 200:1
  - » PE-1(config-vrf)#exit
- Add to the VRF a manual route towars local CE in case of no routing protool on the PE-CE link
  - » PE-1(config)#ip route vrf vpnB 192.168.0.0 255.255.255.0 160.2.11.2
- Associate interface to the VRF
  - » PE-1(config)#int f0/1
  - » PE-1(config-if)#ip vrf forwarding vpnB
  - » PE-1(config-if)#ip address 160.2.11.1 255.255.255.25

- Configure BGP
- Optionally disable IPv4 peering
  - » router bgp 3269
    - » no bgp default ipv4-unicast
- Create peering with all PEs (if not existent)
  - » neighbor 2.2.2.2 remote-as 3269
  - » neighbor 2.2.2.2 update-source Loopback0
  - » neighbor 3.3.3.3 remote-as 3269
  - » neighbor 3.3.3.3 update-source
- Activate vpnv4 peerings
  - » address-family vpnv4
    - » neighbor 2.2.2.2 activate
    - » neighbor 2.2.2.2 send-community extended
    - » neighbor 2.2.2.2 next-hop-self
    - » exit-address-family

- Switch on the BGP advertisements of VRF in case eBGP was not active in the PE-CE link
  - » address-family ipv4 vrf vpnB
  - » network 192.168.0.0

### • In case BGP was active on PE-CE

- » Configure BGP global peering with PE
  - » neighbor 160.2.11.2 remote-as 200
  - » neighbor 160.2.11.2 update-source FastEthernet0/1
- » Bind VRF to the neighbour
  - » address-family ipv4 vrf vpnB
  - » neighbor 160.2.11.2 remote-as 200
  - » neighbor 160.2.11.2 activate
  - » neighbor 160.2.11.2 as-override

### • Debug

- » show ip vrf
- » show ip route vrf vpnB
- » show mpls forwarding-table
  - » Useful to know the external label towards the remote PE
- » show ip bgp vpnv4 vrf vpnB labels
  - » Useful to know the internal label